

Local characterization of decomposability for 2-parameter persistence modules

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PARIS-SACLAY

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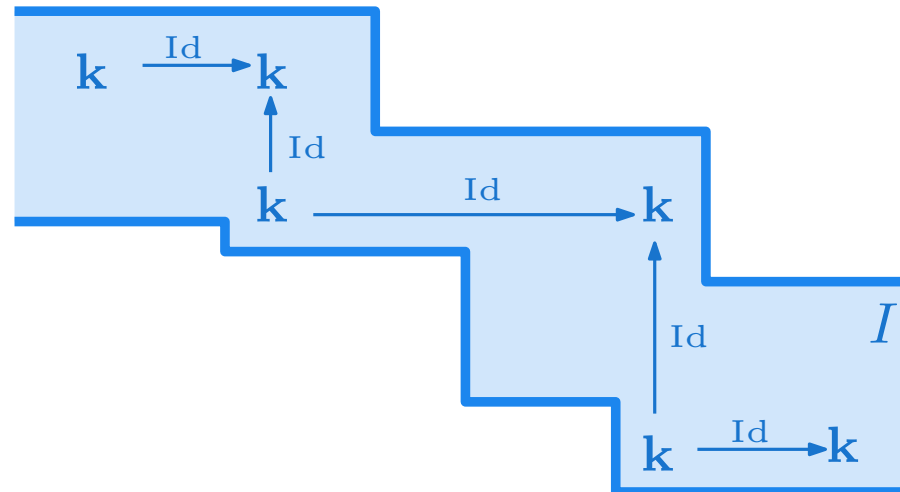
Inria

arxiv:2008.02345

arxiv:2002.08894

joint work with M. Botnan and S. Oudot

(adapted from S. Oudot's slides)



Persistence modules

finite dimensional
k-vector spaces

(P, \leq) a poset, k a field

Persistence module : functor $M : P \rightarrow \text{vec}_k$

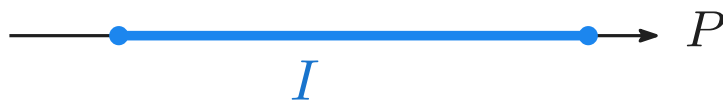
Interval : $I \subseteq P$ that is :

Not. : $I \in \text{Int}(P)$

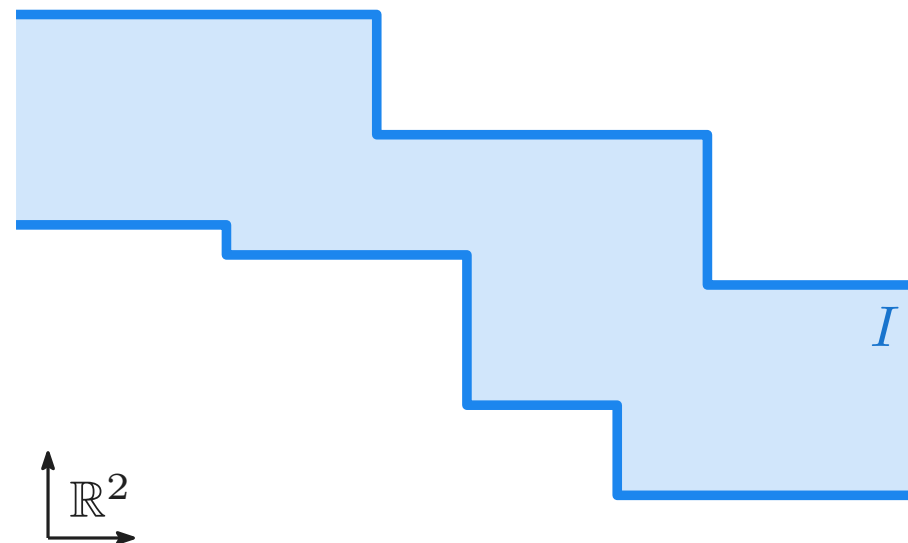
- convex ($s, t \in I \implies u \in I \forall s \leq u \leq t$)

- connected ($s, t \in I \implies \exists \{u_i\}_{i=0}^r \subseteq I$ s.t. $s = u_0 \leq u_1 \geq \dots \geq u_r = t$)

Ex. : (i) P totally ordered



(ii) $P = \mathbb{R}^2$ with coordinatewise order



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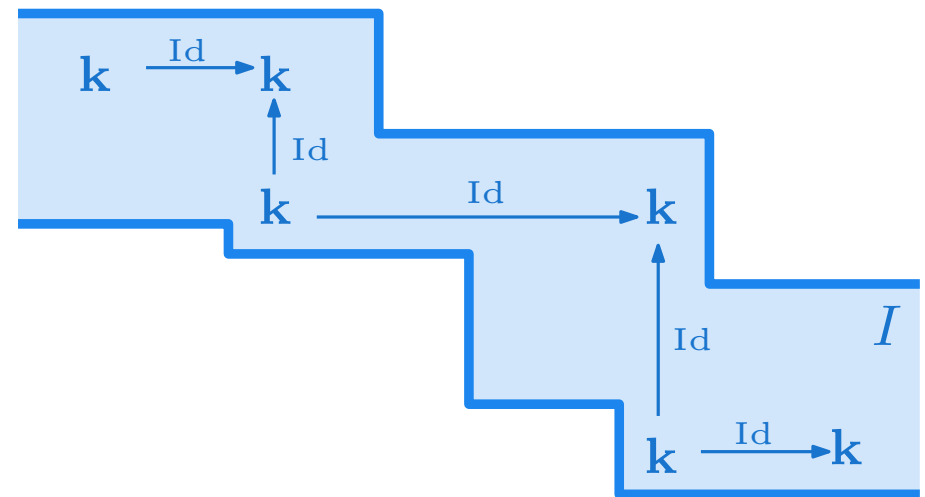
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Interval module : indicator module \mathbf{k}_I of an interval $I \subseteq P$

$$\mathbf{k}_I(t) = \begin{cases} \mathbf{k} & \text{if } t \in I \\ 0 & \text{otherwise} \end{cases}$$

$$\mathbf{k}_I(s \leq t) = \begin{cases} \text{Id}_{\mathbf{k}} & \text{if } s, t \in I \\ 0 & \text{otherwise} \end{cases}$$



Rq : $\text{End}(\mathbf{k}_I) \simeq \mathbf{k} \longrightarrow$ indecomposable

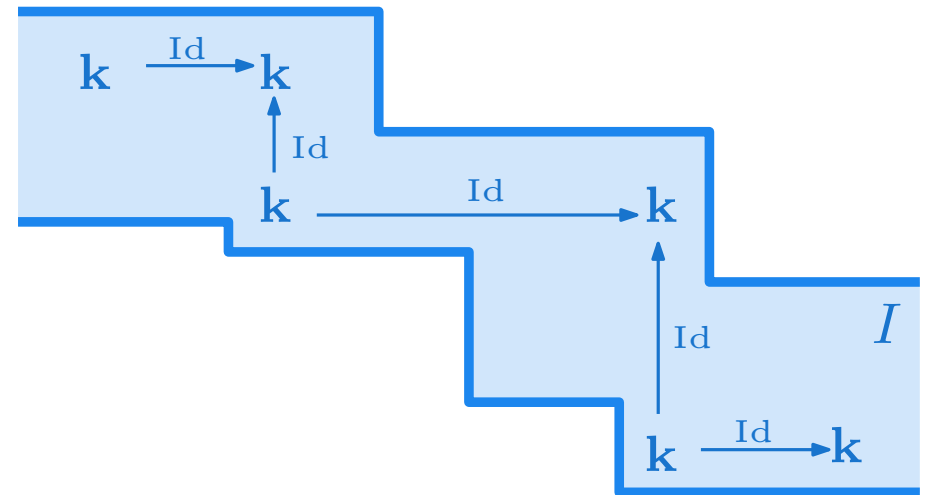
Persistence modules

Interval modules are **described by their support** :

- ▶ geometric descriptor
- ▶ efficient to encode (small / simple dictionary)
- ▶ readily interpretable
- ▶ easy to vectorize (for machine learning)

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Decomposition of persistence modules

Thm. (Crawley-Boevey '15, Botnan, Crawley-Boevey '20)

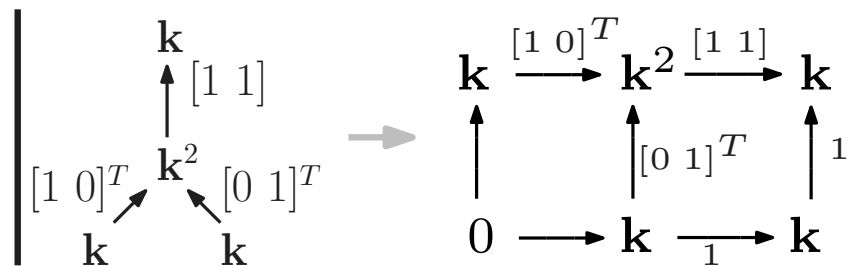
Assume P is totally ordered. Then, $M \simeq \bigoplus_{I \in \mathcal{I}} \mathbf{k}_I$ where the I 's are intervals of P .

Que. What about $P = X_1 \times \cdots \times X_n \subseteq \mathbb{R}^n$ with $n \geq 2$?

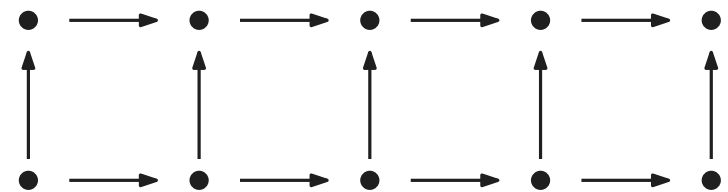
Thm. (Botnan, Crawley-Boevey '20) (P, \leq) any poset

$M \simeq \bigoplus_{\alpha \in A} M_\alpha$ where the M_α 's are indecomposable

Pbs. (i) non-thin indecomposables



(ii) wild type posets



Questions

▶ characterize datasets/filtrations whose modules are interval-dec.

▶ given a filtration : check interval-dec. & extract summands

Local characterizations faster than with a full decomposition (e.g. MeatAxe)

Local characterizations

Poset : $P = X \times Y \subseteq \mathbb{R}^2$

(i) **Collection of supports** : $\mathcal{S} \subseteq \text{Int}(X \times Y)$

Ex. rectangles $\mathcal{S} = \text{Rec}(X \times Y)$
 $\quad = \{ \underline{I} \times \underline{J} : (I, J) \in \text{Int}(X) \times \text{Int}(Y) \}$

associated **\mathcal{S} -decomposable** modules :

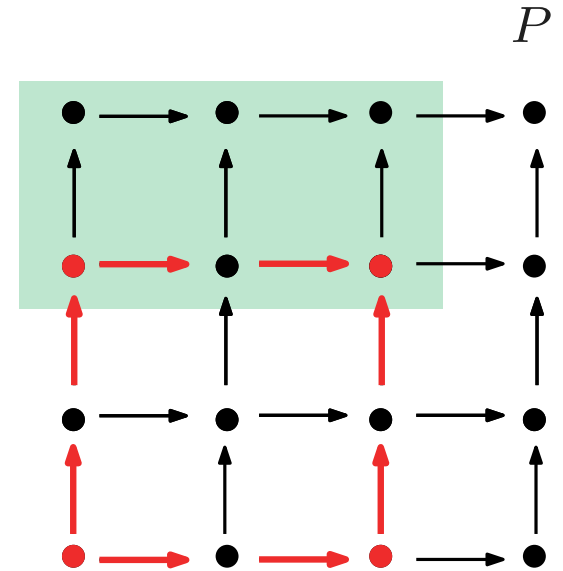
$$\langle \mathcal{S} \rangle = \left\{ M \text{ pers. mod.} : M \simeq \bigoplus_{S \in \mathcal{S}} \mathbf{k}_S^{m_S}, m_S \in \mathbb{Z} \right\}$$

(ii) **Test subsets** : collection \mathcal{Q} of subsets of $X \times Y$

For $Q \in \mathcal{Q}$, $\mathcal{S}|_Q := \{ S \cap Q : S \in \mathcal{S} \} \subseteq \text{Conv}(X \times Y) \longrightarrow \langle \mathcal{S}|_Q \rangle$

Rk. \mathcal{S} -decomposability is closed under taking restrictions, i.e.

$$M \in \langle \mathcal{S} \rangle \implies \forall Q \in \mathcal{Q}, M|_Q \in \langle \mathcal{S}|_Q \rangle$$



Main question

Identify \mathcal{S} and \mathcal{Q} such that : $M \in \langle \mathcal{S} \rangle \iff \forall Q \in \mathcal{Q}, M|_Q \in \langle \mathcal{S}|_Q \rangle$

Known local characterizations

Main question

Identify \mathcal{S} and \mathcal{Q} such that : $M \in \langle \mathcal{S} \rangle \iff \forall Q \in \mathcal{Q}, M|_Q \in \langle \mathcal{S}|_Q \rangle$

Setting : $\mathcal{S} = \text{Blc}(X \times Y)$ $\mathcal{Q} = \{\text{squares of } X \times Y\}$

$$\text{Blc}(X \times Y) = \left\{ \begin{array}{ll} \begin{array}{|c|} \hline \square \\ \hline \end{array} & \text{"death quadrants"} \\ \begin{array}{|c|} \hline \square \\ \hline \end{array} & \text{"birth quadrants"} \end{array} \right\} \quad \left\{ \begin{array}{ll} \begin{array}{|c|} \hline \square \\ \hline \end{array} & \text{"horizontal bands"} \\ \begin{array}{|c|} \hline \square \\ \hline \end{array} & \text{"vertical bands"} \end{array} \right\}$$

Rk. For any square Q ,

$$\text{Rec}(X \times Y)|_Q = \left\{ \begin{array}{cccccccc} \bullet & \bullet & \bullet & \bullet & \text{---} & \text{---} & \bullet & \bullet \\ \bullet & \bullet & \bullet & \bullet & \bullet & \bullet & \bullet & \bullet \end{array} \right\}$$

$$\text{Blc}(X \times Y)|_Q = \left\{ \begin{array}{cccccccc} \bullet & \bullet & \bullet & \bullet & \text{---} & \text{---} & \bullet & \bullet \\ \bullet & \bullet & \bullet & \bullet & \bullet & \bullet & \bullet & \bullet \end{array} \right\}$$

Thm. (Cochoy, Oudot '20) $\mathcal{S} = \text{Blc}(X \times Y)$ $\mathcal{Q} = \{\text{squares of } X \times Y\}$

M is block-dec. iff its restriction to any square is block-dec., i.e.

$$M \in \langle \text{Blc}(X \times Y) \rangle \iff \forall Q \in \mathcal{Q}, M|_Q \in \langle \text{Blc}(Q) \rangle$$

Silly question

Case : $\mathcal{S} = \text{Int}(X \times Y)$ and $\mathcal{Q} = \{\text{totally ordered subsets of } X \times Y\}$

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Fact. Interval-dec. cannot be checked on totally ordered subsets.

Proof. restrictions to totally ordered subsets are **always** interval-dec (Crawley-Boevey).

$$M := \begin{array}{ccccc} & & \begin{matrix} [1] \\ 0 \end{matrix} & & \\ & & \longrightarrow & & \\ \mathbf{k} & \xrightarrow{\quad} & \mathbf{k}^2 & \xrightarrow{\quad [1 \ 1] \quad} & \mathbf{k} \\ \uparrow & & \uparrow & & \uparrow \\ 0 & & \begin{matrix} [0] \\ 1 \end{matrix} & & 1 \\ \uparrow & & \uparrow & & \uparrow \\ 0 & \xrightarrow{\quad 0 \quad} & \mathbf{k} & \xrightarrow{\quad 1 \quad} & \mathbf{k} \end{array}$$

indecomposable

not an interval module

Interval-decomposability is not local

Case : $\mathcal{S} = \text{Int}(X \times Y)$ and $\mathcal{Q} = \{\text{finite strict subgrids of } X \times Y\}$

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Thm. (Botnan, L., Oudot '20)

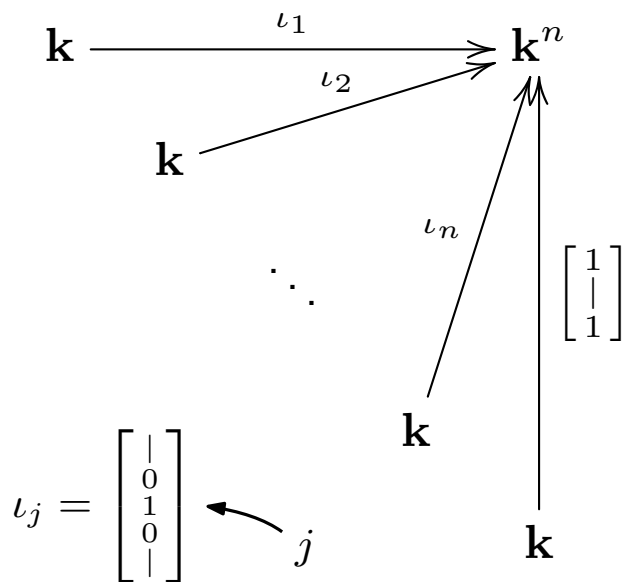
Assume that $\#X > 2$ and $\#Y > 2$. Let $2 \leq m < \min(\#X, \#Y)$.

There exists a module M over $X \times Y$ such that :

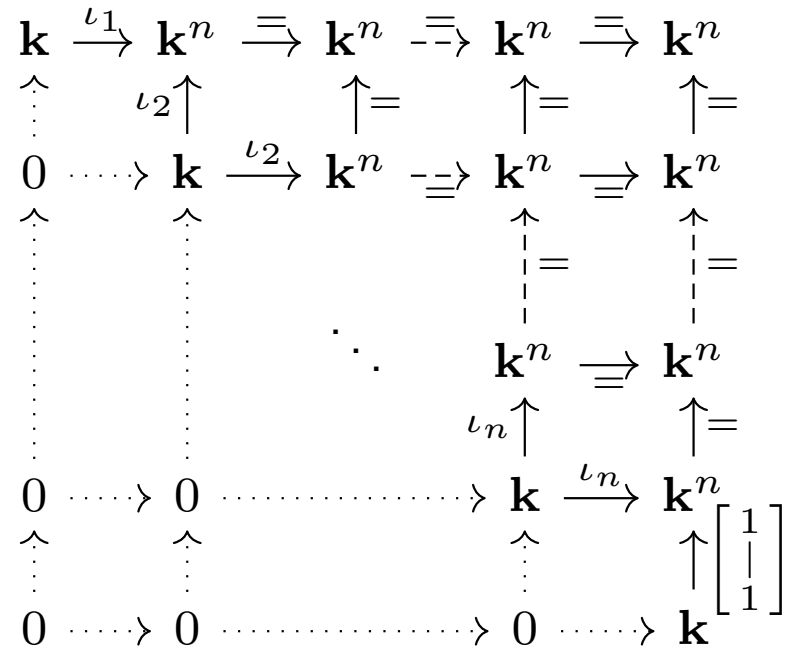
- (i) M is not interval-dec.
- (ii) $M|_Q$ is interval-dec for any grid $Q \subset X \times Y$ of side-length at most m .

Proof :

Adapt the indecomposable :



$\rightarrow M :=$
 indec.
 not interval module



Rectangle-decomposability is local

Case : $\mathcal{S} = \text{Rec}(X \times Y)$ and $\mathcal{Q} = \{\text{squares of } X \times Y\}$

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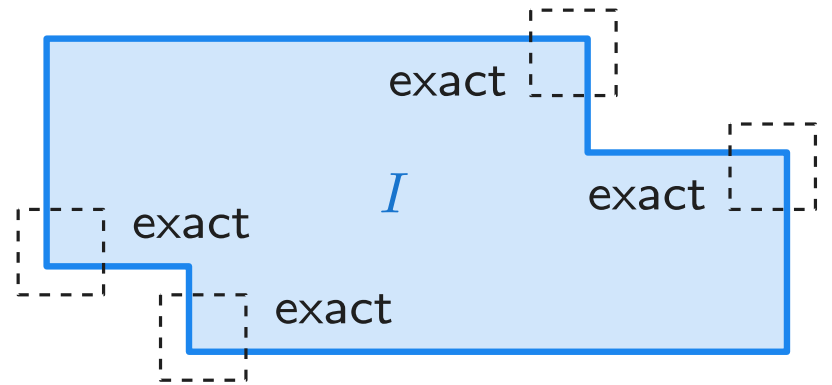
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Algebraic characterization

$$\begin{array}{ccc} M(c) & \xrightarrow{\delta} & M(d) \\ \beta \uparrow & & \uparrow \gamma \\ M(a) & \xrightarrow{\alpha} & M(b) \end{array}$$



Exactness condition : $M(a) \xrightarrow{\varphi = (\alpha, \beta)} M(b) \oplus M(c) \xrightarrow{\psi = \gamma - \delta} M(d)$ is exact.

Ex. Interlevel sets persistent homology

$$(x, y) \in \mathbb{R}^{\text{op}} \times \mathbb{R} \mapsto H_k(f^{-1}(x, y))$$

Ex.

$$\begin{array}{ccc} & \text{exact} & \\ & 0 \longrightarrow 0 & \\ \uparrow & & \uparrow \\ \mathbf{k} & \longrightarrow 0 & \end{array}$$

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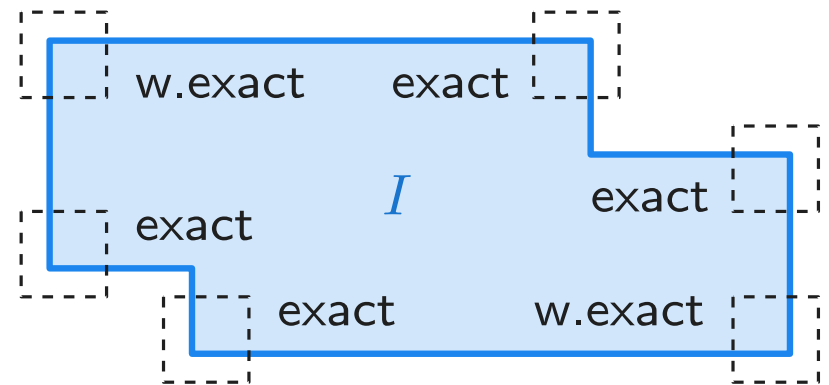
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Weak exactness condition :

$$\begin{array}{|l}
 \text{Im } \rho = \text{Im } \gamma \cap \text{Im } \delta \\
 \text{Ker } \rho = \text{Ker } \alpha + \text{Ker } \beta
 \end{array}$$

Ex.

exact	w.exact	not w.exact
$0 \longrightarrow 0$	$\mathbf{k} \longrightarrow 0$	$\mathbf{k} \longrightarrow 0$
$\uparrow \qquad \uparrow$	$\uparrow \qquad \uparrow$	$\uparrow \qquad \uparrow$
$\mathbf{k} \longrightarrow 0$	$0 \longrightarrow 0$	$\mathbf{k} \longrightarrow \mathbf{k}$

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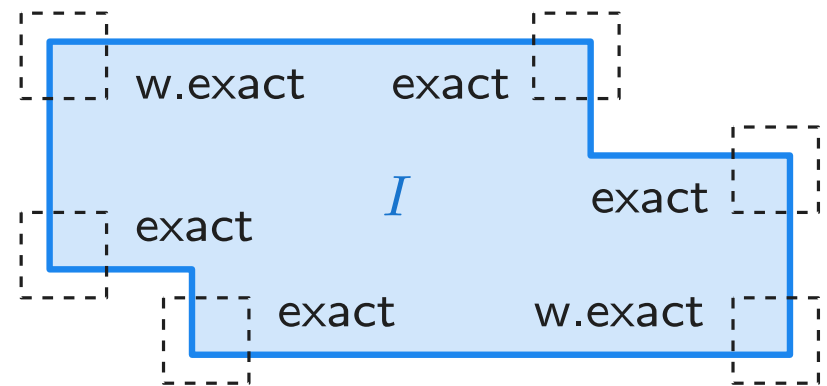
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Lem.

$$M \text{ is exact} \iff M|_Q \in \langle \text{Blc}(Q) \rangle$$

$$M \text{ is w. exact} \iff M|_Q \in \langle \text{Rec}(Q) \rangle$$

for all
squares
 Q

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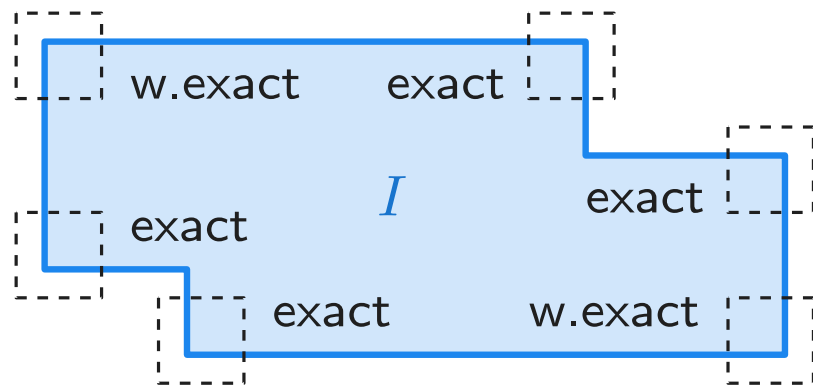
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$$M \in \langle \text{Rec}(X \times Y) \rangle \iff \underline{M \text{ is weakly exact}}$$

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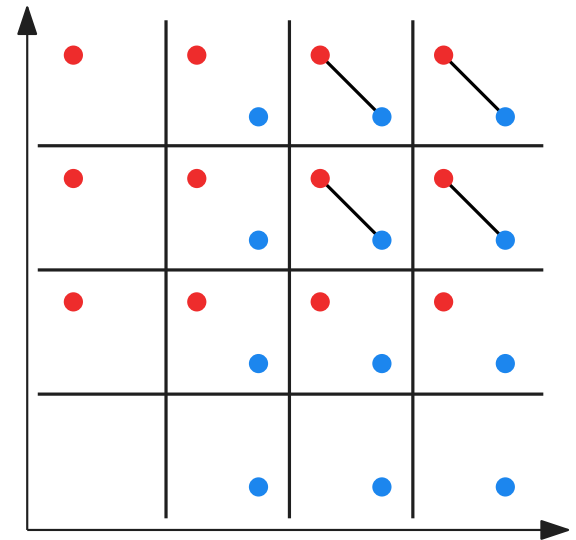
Application : checking rectangle-dec.

$$P = \{1, \dots, n\} \times \{1, \dots, m\} \subset \mathbb{R}^2$$

Given $F : P \rightarrow \text{Simp}$ (finite simplicial complexes) :

Goal :

- ▶ determine whether $H_k F$ is rectangle-dec.
- ▶ if so, compute the decomposition of $H_k F$



Straightforward approach

- compute direct-sum decomposition $\longrightarrow O(r^{2\omega+1})$ (Dey, Xin '19)
- check summands one by one (thin-ness and support)

Time complexity : $O(r^{2\omega+1})$

r : total number of simplices in F

$\omega \approx 2.373$: exponent for matrix multiplication

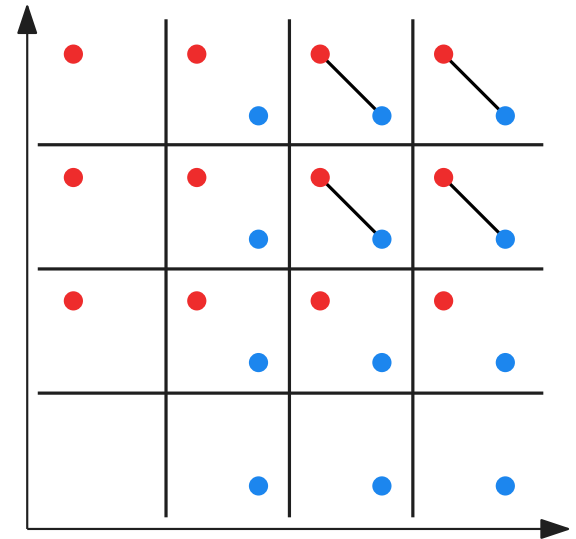
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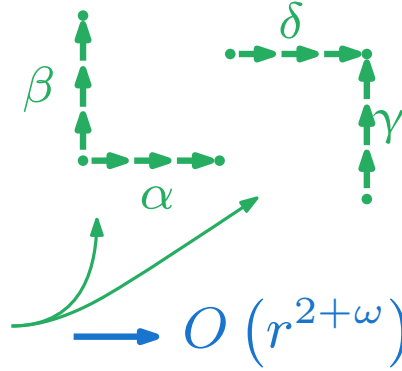
Optimized approach for rectangle-dec.

(i) compute rank invariant $\text{rk} : P \times P \rightarrow \mathbb{N} \rightarrow O(r^4)$

(ii) compute all $\dim(\text{Im } \gamma \cap \text{Im } \delta)$ and $\dim(\text{Ker } \alpha + \text{Ker } \beta)$

Method : compute $2 \times r^2$ zigzag barcodes in $O(r^\omega)$ each $\rightarrow O(r^{2+\omega})$

(iii) check weak exactness, i.e. $\dim \text{Im } \rho = \dim(\text{Im } \gamma \cap \text{Im } \delta)$
 $\dim \text{Ker } \rho = \dim(\text{Ker } \alpha + \text{Ker } \beta) \rightarrow O(r^4)$



Time complexity : $O(r^{2+\omega})$

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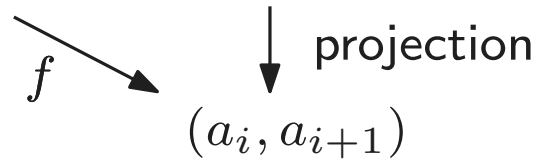
$\omega \approx 2.373$: exponent for matrix multiplication

Application : Pyramid basis theorem

Setting : $f : X \rightarrow \mathbb{R}$ of “Morse type”

i.e. with “critical values” $a_0 = -\infty < a_1 < \dots < a_n < a_{n+1} = +\infty$

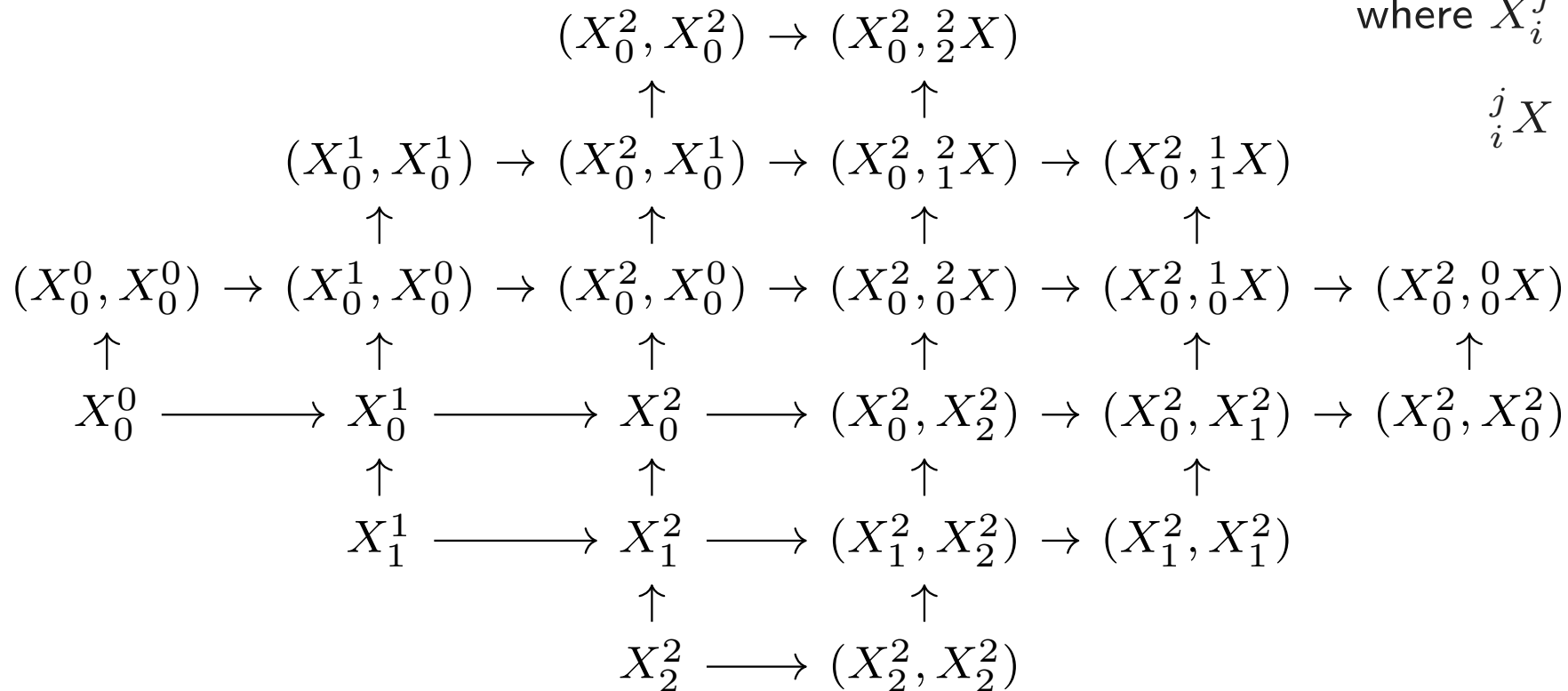
s.t. $f^{-1}(a_i, a_{i+1}) \simeq Y_i \times (a_i, a_{i+1})$



Consider $-\infty < s_0 < a_1 < \dots < a_n < s_n < +\infty$. We have inclusions :

where $X_i^j := f^{-1}(s_i, s_j)$

${}^j_i X := f^{-1}(s_0, s_i) \cup f^{-1}(s_j, s_n)$

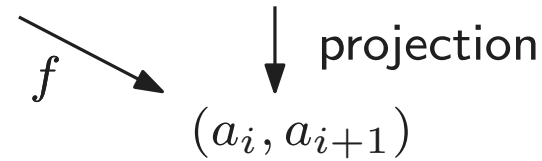


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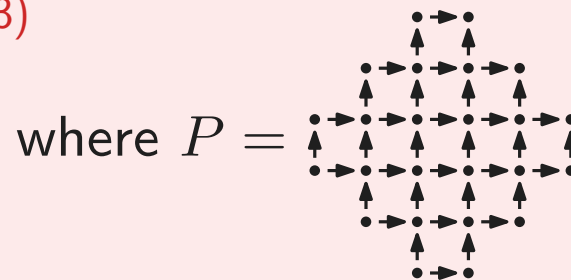
Taking (relative) homology :

$$\begin{array}{cccccccc}
 & & & & H_p(X_0^2, X_0^2) & \longrightarrow & H_p(X_0^2, {}^2_2 X) & & \\
 & & & & \uparrow & & \uparrow & & \\
 & & & & H_p(X_0^1, X_0^1) & \longrightarrow & H_p(X_0^2, X_0^1) & \longrightarrow & H_p(X_0^2, {}^2_1 X) & \longrightarrow & H_p(X_0^2, {}^1_1 X) & & \\
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 H_p \mathcal{P} := & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \\
 & & & & H_p(X_0^0) & \longrightarrow & H_p(X_0^1) & \longrightarrow & H_p(X_0^2) & \longrightarrow & H_p(X_0^2, X_2^2) & \longrightarrow & H_p(X_0^2, X_1^2) & \longrightarrow & H_p(X_0^2, X_0^2) & & \\
 & & & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \\
 & & & & H_p(X_1^1) & \longrightarrow & H_p(X_1^2) & \longrightarrow & H_p(X_1^2, X_2^2) & \longrightarrow & H_p(X_1^2, X_1^2) & & & & & & \\
 & & & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & & & & & \\
 & & & & H_p(X_2^2) & \longrightarrow & H_p(X_2^2, X_2^2) & & & & & & & & & &
 \end{array}$$

Application : Pyramid basis theorem

Thm. (Bendich, Edelsbrunner, Morozov, Patel, '13)

$$H_p \mathcal{P} \simeq \bigoplus_{\substack{R \in \text{Rec}(\mathbb{Z}^2) \\ \text{s.t. } R \cap P \text{ is maximal in } P}} \mathbf{k}_{R \cap P}^m$$



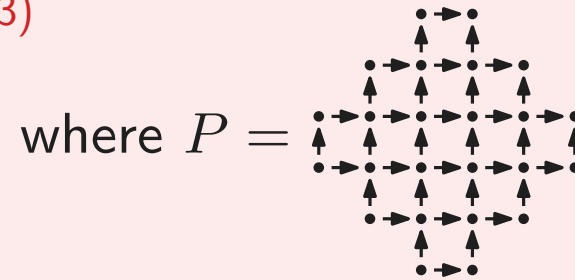
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 & & H_p(X_0^0) & \longrightarrow & H_p(X_0^1) & \longrightarrow & H_p(X_0^2) & \longrightarrow & H_p(X_0^2, X_2^2) & \longrightarrow & H_p(X_0^2, X_1^2) & \longrightarrow & H_p(X_0^2, X_0^2) \\
 & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow \\
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 & & \uparrow & & \uparrow & & \uparrow & & \uparrow \\
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 \end{array}$$

Application : Pyramid basis theorem

Thm. (Bendich, Edelsbrunner, Morozov, Patel, '13)

$$H_p \mathcal{P} \simeq \bigoplus_{\substack{R \in \text{Rec}(\mathbb{Z}^2) \\ \text{s.t. } R \cap P \text{ is maximal in } P}} \mathbf{k}_{R \cap P}^{m_R}$$



New proof. (i) $\text{Rk} : H_p \mathcal{P}$ is **exact**.

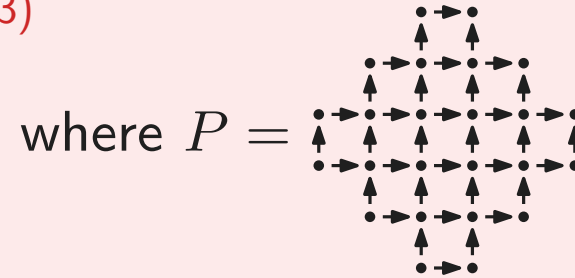
(ii) Extend $H_p \mathcal{P}$ to a product $X \times Y$, into a **weakly exact** module :

$$\begin{array}{cccccccc}
 & & & & H_p(X_0^2, X_0^2) & \longrightarrow & H_p(X_0^2, {}_2X) & & \\
 & & & & \uparrow & & \uparrow & & \\
 & & & & H_p(X_0^1, X_0^1) & \longrightarrow & H_p(X_0^2, X_0^1) & \longrightarrow & H_p(X_0^2, {}_2X) & \longrightarrow & H_p(X_0^2, {}_1X) & & \\
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 & & & & H_p(X_0^0) & \longrightarrow & H_p(X_0^1) & \longrightarrow & H_p(X_0^2) & \longrightarrow & H_p(X_0^2, X_2^2) & \longrightarrow & H_p(X_0^2, X_1^2) & \longrightarrow & H_p(X_0^2, X_0^2) & & \\
 & & & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \\
 & & & & H_p(X_1^1) & \longrightarrow & H_p(X_1^2) & \longrightarrow & H_p(X_1^2, X_2^2) & \longrightarrow & H_p(X_1^2, X_1^2) & & & & & & \\
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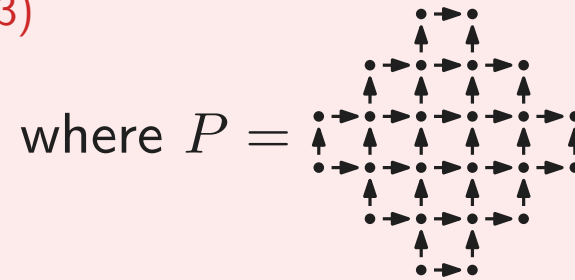
$$\begin{array}{cccccccccccc}
 0 & \longrightarrow & 0 & \longrightarrow & H_p(X_0^2, X_0^2) & \longrightarrow & H_p(X_0^2, \frac{2}{2}X) & \longrightarrow & \text{PO}_2 & \longrightarrow & \text{PO}_3 \\
 \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow \\
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 \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow \\
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$H_p \mathcal{P} :=$

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 \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow \\
 \text{PB}_1 & \longrightarrow & H_p(X_1^1) & \longrightarrow & H_p(X_1^2) & \longrightarrow & H_p(X_1^2, X_2^2) & \longrightarrow & H_p(X_1^2, X_1^2) & \longrightarrow & 0 \\
 \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow & & \uparrow \\
 \text{PB}_3 & \longrightarrow & \text{PB}_2 & \longrightarrow & H_p(X_2^2) & \longrightarrow & H_p(X_2^2, X_2^2) & \longrightarrow & 0 & \longrightarrow & 0
 \end{array}$$

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Rectangle-dec. is the biggest local class

Case : $\mathcal{S} \subseteq \text{Int}(X \times Y)$ s.t. for all square Q in $X \times Y$:

$$\mathcal{S}|_Q \supseteq \left\{ \begin{array}{cccccccccccc} \bullet \bullet & ; & \bullet \bullet & ; & \bullet \bullet & ; & \bullet \bullet & ; & \text{---} & ; & \text{---} & ; & \text{I} \bullet & ; & \bullet \text{I} & ; & \blacksquare & ; & \text{L} \end{array} \right\}$$

rectangles
"hook"

$$Q = \{\text{squares of } X \times Y\}$$

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Thm. (Botnan, L., Oudot '20)

Assume that $\#X \geq 2$ and $\#Y \geq 2$ but $(\#X, \#Y) \neq (2, 2)$.

There exists a module M over $X \times Y$ such that :

(i) M is not $\langle \mathcal{S} \rangle$ -dec.

(ii) $M|_Q$ is $\langle \mathcal{S}|_Q \rangle$ -dec for any square Q of $X \times Y$.

Proof :

$$\begin{array}{ccccc}
 \mathbf{k} & \xrightarrow{\begin{bmatrix} 1 \\ 0 \end{bmatrix}} & \mathbf{k}^2 & \xrightarrow{\begin{bmatrix} 1 & 1 \end{bmatrix}} & \mathbf{k} \\
 \uparrow 0 & & \uparrow \begin{bmatrix} 0 \\ 1 \end{bmatrix} & & \uparrow 1 \\
 0 & \xrightarrow{0} & \mathbf{k} & \xrightarrow{1} & \mathbf{k}
 \end{array}$$

\longrightarrow indecomposable not interval module

\longrightarrow $\langle \mathcal{S}|_Q \rangle$ -decomposable for all square Q

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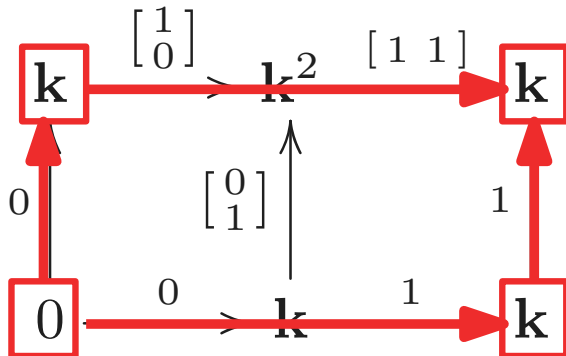
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→ indecomposable not interval module

→ $\langle \mathcal{S}|_Q \rangle$ -decomposable for all square Q

Ideas of proof : interval-decomposition (Crawley-Boevey)

1. Define a *counting functor* for each interval I :

$$C_I : \text{vec}_{\mathbf{k}}^{\mathbb{R}} \rightarrow \text{vec}_{\mathbf{k}}$$

$$M \mapsto \mathbf{k}^{\text{mult}(\mathbf{k}_I; M)} \quad \text{where } \text{mult}(\mathbf{k}_I; M) := \max\{n \mid M \simeq \mathbf{k}_I^n \oplus N\}$$

2. Define a submodule M_I of M for each interval I such that :

$$M_I \simeq \mathbf{k}_I^{\text{mult}(\mathbf{k}_I; M)}$$

3. Show that $M = \bigoplus_I M_I$

- show that the M_I 's are in direct sum
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Ideas of proof : interval-decomposition

(Crawley-Boevey)

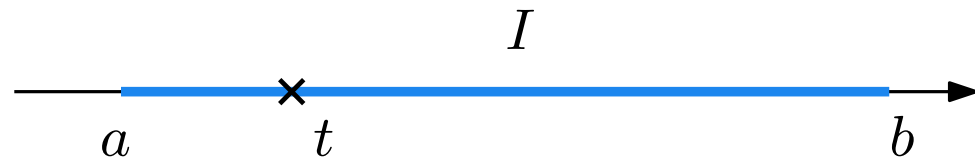
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Details : For $I = (a, b)$:



- $\text{Im}_I^+(t) := \bigcap_{a < s \leq t} \text{Im } M(s \rightarrow t)$

(elements alive at least since a and still at t)

- $\text{Im}_I^-(t) := \sum_{s \leq a} \text{Im } M(s \rightarrow t)$

(elements born before a and still alive at t)

↳ $\text{Im}_I^+(t) / \text{Im}_I^-(t)$

(elements alive at t that were born at a)

- $\text{Ker}_I^+(t) := \bigcap_{s \geq b} \text{Ker } M(t \rightarrow s)$

(elements alive at t but not after b)

- $\text{Ker}_I^-(t) := \sum_{t \leq s < b} \text{Ker } M(s \rightarrow t)$

(elements alive at t and dead before b)

↳ $\text{Ker}_I^+(t) / \text{Ker}_I^-(t)$

(elements alive at t that die at b)

Ideas of proof : interval-decomposition

(Crawley-Boevey)

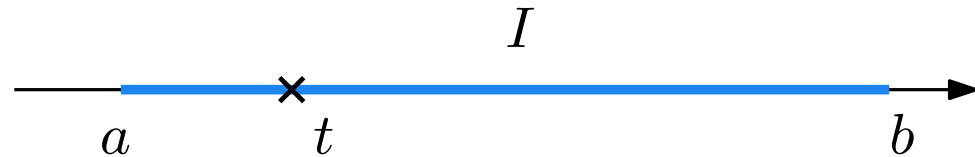
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Details : For $I = (a, b)$:



$$C_I[t] := \frac{\text{Im}_I^+(t) \cap \text{Ker}_I^+(t)}{\left(\text{Im}_I^+(t) \cap \text{Ker}_I^-(t) \right) + \left(\text{Im}_I^-(t) \cap \text{Ker}_I^+(t) \right)}$$

(alive at least since a but not after b)



(alive since a but dead before b) + (alive until b but born before a)

Prop. For $t \leq t' \in (a, b)$, $M(t \rightarrow t')$ induces $C_I(t) \xrightarrow{\simeq} C_I(t')$

Def. $C_I(M) := \varprojlim_{t \in I} C_I(t)$

Ideas of proof : interval-decomposition (Crawley-Boevey)

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Idea of proof : rectangle-decomposition

1. Define a *counting functor* for each **rectangle** R :

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2. Define a submodule M_R of M for each **rectangle** R such that :

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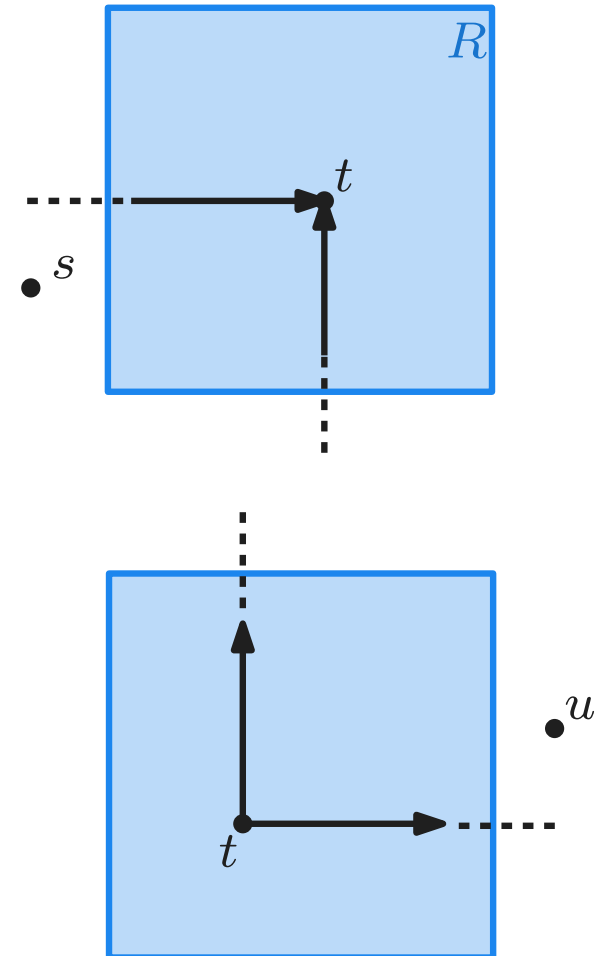
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Pb : product order on \mathbb{R}^2 is not total

$$\sum_{\substack{s \notin R \\ s \leq t}} \text{Im } M(s \rightarrow t) \not\subseteq \bigcap_{\substack{s \in R \\ s \leq t}} \text{Im } M(s \rightarrow t)$$

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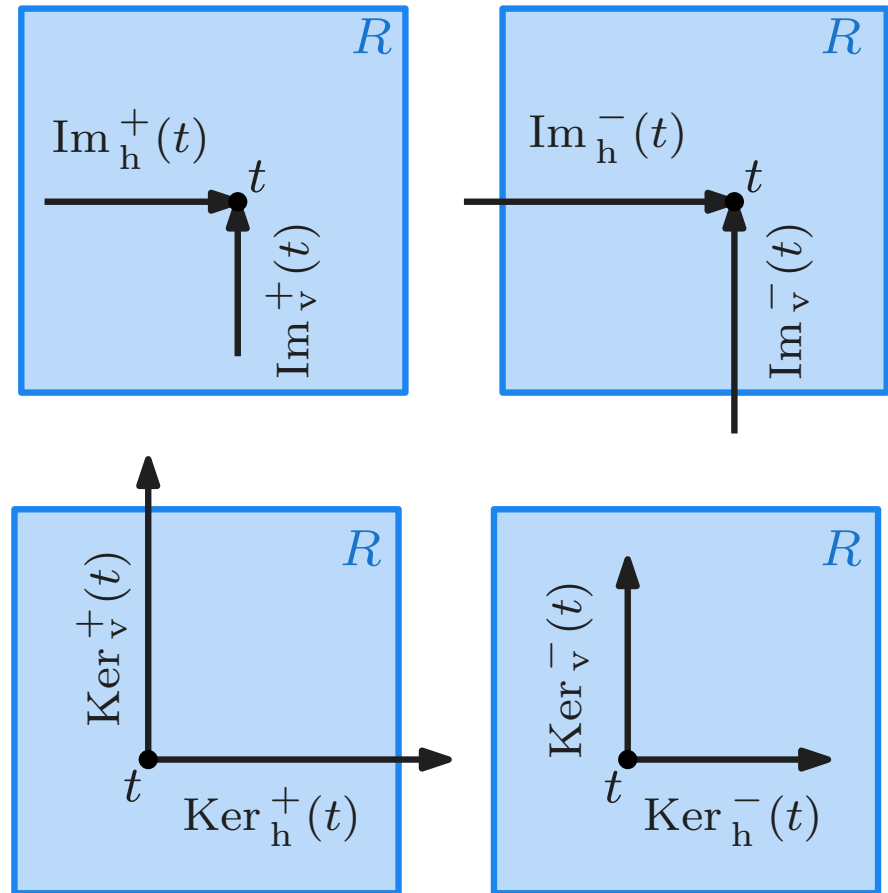
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Sol. : decouple hori./vert. contributions using **weak exactness**

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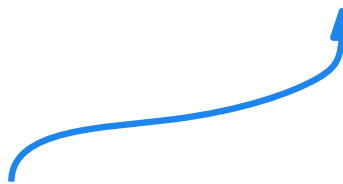
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Then : Define $C_R(M)$ similarly



$$\text{Im}_R^-(t) := \left(\text{Im}_h^-(t) + \text{Im}_v^-(t) \right) \cap \text{Im}_R^+(t)$$

$$\text{Ker}_R^+(t) := \text{Ker}_R^-(t) \\ + \left(\text{Ker}_h^+(t) \cap \text{Ker}_v^+(t) \right)$$

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Conclusion

Summary

- (i) Interval decomposability is not local
- (ii) Block & rectangle decomposability are local
- (iii) Rectangle decomposability is the biggest local subclass of interval-dec.

Open questions

Partial local characterization ?

- ▶ existence of interval summands in the decomposition
- ▶ extraction of those interval summands

Classes \mathcal{I} of indecomposables beyond $\text{Int}(P)$?

- ▶ local characterization
- ▶ compute decomposition

Posets beyond \mathbb{R}^2 ?

- ▶ \mathbb{R}^d
- ▶ others ?

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Thank you !

arxiv:2008.02345

arxiv:2002.08894